M-MENTOR: A Design Algorithm for IP Networks with Mixed Traffic

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Abstract: The design of IP networks to support traffic engineering for both unicast and multicast traffic is a very difficult problem. This paper proposes a heuristic design algorithm called M-MENTOR that concern routing of both types of traffic. However, since multicast traffic model could be employed in many situations and could be managed by various kinds of multicast routing protocols, this paper consider only the design process of the IP networks with following features: (1) network of within the same Autonomous System (AS), (2) routing protocol support multiple weight on each link, (3) the multicast traffic from different sources share the same multicast tree. The efficiency of M-MENTOR is evaluated for various traffic demands and networks of 10, 30 and 50 node and compared with MENTOR-II. The experimental results show that, in almost all cases, M-MENTOR give better performance in term of installation cost.

Key-Words: - IP Network Design, MENTOR Algorithm, Unicast/Multicast Traffic, Traffic Engineering

1 Introduction

IP network design that concern both unicast and multicast routing is a very difficult problem. The problem is even more difficult if we chose mange the traffic by appropriate weight setting of links in the OSPF protocol instead of using the overlay network technique. These kinds of problems are classified as "Mixed Integer Linear Programming" or MIP [1].

To reduce complexity of the network design process Kershenbaum et al. [2] has developed a complexity heuristic algorithm with low called MENTOR (Mesh Network Topological Optimization and Routing). The networks designed by MENTOR give very good performance that very close to that of the optimum [3][4]. MENTOR also can be used to design virtual circuits packet switching network such as ATM or Frame Relay. However, a MENTOR algorithm can not be used directly to design routers or MPLS routers networks that employ OSPF or IS-IS routing protocol. This is because MENTOR does not perform appropriate link weight setting. Cahn [5] has improved the MENTOR algorithm such that appropriate OSPF link weight can be set during the design process using Incremental Shortest Path (ISP). Such algorithm is known as the MENTOR-II. However, it should be note that almost all the above design algorithms have been developed for network with only unicast traffic.

Presently, several important emerging multicast applications such as distributed database systems, radio, television, video conferencing system, distance learning system, are becoming more and more popular. As a result, the portion of multicast traffic on the IP network in almost all organization is increasing rapidly. Therefore, IP network design process should also effectively route multicast traffic in addition to the traditional unicast traffic.

The efficient IP networks design processes for multicast traffic should to consider the traffic traversing paths and routing protocols. For multicast traffic, it is well known that the optimum spanning tree that connects the transmitter and all its receivers is called Steiner Tree. Unfortunately, there is no known algorithm that can be used to systematically construct the Steiner Tree. Today, most Steiner trees are obtained by heuristic algorithm. One of the best known heuristic for Steiner is T-M algorithm that was proposed by Takahashi and Matsuyama [6]. An interesting modification of T-M algorithm has been proposed by Maxemchuk [10] which is optimized for multi speed multicast systems.

There are many choices for multicast routing protocols. But, one of the most popular used within an Autonomous System (AS) is Protocol Independent Multicast (PIM). PIM also has several modes depend upon the density or the number of receivers within the area. PIM-Dense Mode (DM) is for the case of high receiver density. PIM-Any Source Multicast (ASM) is for the case of low receiver density and there are few numbers of transmitters, e.g. few-to-many. PIM-Source Specific Multicast (SSM) is for the case single transmitter, e.g. one-to-many. PIM-Bider (Bidirectional) is for the case that there are many receivers and transmitters (Many-to-many). PIM-DM is now obsolete. The most popular option may be PIM-ASM because it supports both PIM-DM and PIM-SSM modes. However, PIM-Bider is seemed to be the most efficient because transmitters share the same multicast tree. With shared multicast trees, the routers do not have to remember different tree for each transmitter.

The problem of PIM is that, it uses the link weight based Shortest Path Tree to distribute the traffic. Theoretically, this Shortest Path Trees are not necessary Steiner Trees which are more favorable for multicast operation.

To improve routing performance for mixed unicast and multicast traffic Wang and Pavlou [7] proposed to separate the flow management of unicast and multicast traffic by employing a newly proposed Interior Gateway Routing Protocol (IGP) that support multiple sets of link weights called Multitopology Enable IGP (MT-IGP) Multitopology extension to ISIS (M-ISIS) [8] and Multitopology extension to OSPF (M-OSPF) [9]. However, Wang and Pavlou [7] have not yet discussed about the network design process for mixed unicast and multicast traffic.

This study proposes a heuristic design algorithm for IP networks with mixed unicast and multicast traffic called M-MENTOR. The algorithm is obtained by modifying spanning tree building portion of MENTOR-II. M-MENTOR used modified T-M algorithm to construct core spanning tree instead of Prim-Dijkstra algorithm. Therefore, M-MENTOR is focus on design of IP networks with following features: (1) all network members are within the same AS, (2) the employ routing protocol should support multiple link weight such as M-OSPF and M-ISIS, (3) the multicast traffic from different sources share the same multicast tree. e.g. PIM-Bidir is deployed. The performances of networks design by M-MENTOR are evaluated in term of installation cost and compared with the networks design by original MENTOR-II for various traffic demands and networks with different number of nodes. The experimental results show that, in almost all cases, M-MENTOR gives better performance in term of installation cost.

2 Backgrounds

2.1 MENTOR Algorithm

MENTOR algorithm [2] is a low complexity heuristic network design algorithm. This low complexity is achieved by doing implicit routing over a link at the same time it is considered to be installed. For a given set of nodes N, demand matrix D and link cost matrix X, let $d_{s,t}$ and $x_{s,t}$ are the amount of traffic flow and link installation cost from s and t, respectively. The properties of network obtained by MENTOR algorithm are (1) traffic demands are routed on relatively direct paths (2) links have reasonable utilization and (3) relatively high capacity links are used. MENTOR starts with clustering process. In this stage, nodes are classified in to end nodes and backbone nodes using a clustering algorithm. Examples of possible clustering algorithms are threshold clustering and Kmean clustering. Here in this paper, we consider only the case where traffic demands are distributed equivalently among all nodes. Therefore, all nodes can be considered as backbone node.

Next, a good tree is formed to interconnect all (backbone) nodes. Kershenbaum et. al. [2] suggests to a use a heuristic, which can be thought of as a modification of Prim and Dijkstra algorithm to build the tree. The algorithm works almost the same manner as Dijkstra algorithm but with a tunable parameter α , $0 \le \alpha \le 1$. The tree is to be expanded one node at a time by connecting a tree node *i* to an out of tree node *j* such that $\alpha L_i + x_{ic}$ minimized, where L_i is the cost of path from root node along the tree to node *i*. Note that $\alpha = 0$ and 1 is corresponding to Minimum Spanning Tree (MST) and Shortest Path Tree (SPT), respectively

Given a tree, the objective of MENTOR is to adding a direct link between each pair of nodes if the amount of traffic is reasonable. Let the maximum utilization be ρ , and the minimum utilization be defined in term of ρ and slack s as $(1-s) \rho$, where s, $0 \le s \le 1$. Consider a pair of nodes A and B, let C_{AB} and l_{AB} be link capacity and accumulated load flow between A and B, respectively. If traffic between A and B is too small, i.e. $l_{AB} < \rho C_{AB}$ (1s), no link is added and all traffic l_{AB} is overflowed to the next most direct path. A link is added if traffic is in between maximum and minimum utilization, i.e. ρC_{AB} $(1-s) \le l_{AB} \le \rho C_{AB}$. However, if $l_{AB} > \rho C_{AB}$, a direct link is added only when traffic bifurcation among multiple routes is possible. If bifurcation is possible, a new link of C_{AB} is added to serve a portion of traffic ρ C_{AB} , and the left portion $l_{AB} - \rho C_{AB}$ is overflowed to the next most direct path. Otherwise, if no bifurcation is possible, no link is added and all traffic l_{AB} is overflowed to the next most direct path.

2.2 MENTOR II Algorithm

MENTOR-II [2] also starts with node clustering and building a good spanning tree between backbone nodes. But when MENTOR-II consider adding a direct link to serve traffic demand between a pair of nodes, at the same time, calculates an appropriate link weight for the each link based on Incremental Shortest Path (ISP) algorithm. The concepts of MENTOR-II and ISP algorithm can be described as follows:

1) At start, set the weight for each link in the spanning tree to the installation cost of the link.

- 2) Let $d_{spt}(A,B)$ be the shortest path distance between node A and B through the selected good spanning tree, consider adding a direct link between each nodes pair in the decreasing order of $d_{spt}(\cdot)$.
- 3) When consider whether to add a link L_{AB} between A and B, the weight w_{AB} of L_{AB} is initially set to a reasonably high value. ISP then tries to draw traffic flow through L_{AB} as much as possible by lowering the w_{AB} . A constraint is that the link w_{AB} should be greater or equal to the installation cost.
- 4) L_{AB} is added if we can find an eligible value of w_{AB} and the amount of traffic flow though it falls in the reasonable zone defined by ρ , C_{AB} , and s.

When MENTOR-II considers all possible direct links, all links are assigned with appropriate link weights which ensure the shortest path routing.

2.3 The Steiner Tree and T-M Heuristic

Given a directional graph G(E, V), let $S \subseteq V$ be a set of nodes involved in multicast communication. Our objective is to construct a minimum cost tree interconnecting all members of S. The member of the tree may include some node $n \notin V$. but $n \in S$. In this study, the cost of a link is an increasing function of the bandwidth and the length of the link, and not dependent upon the utilization or availability of the link. Such problem is called Steiner tree problem. A special case is when $S \equiv V$, Steiner tree is equivalent to Minimum Spanning Tree (MST). However, unlike MST problem, there is no systematic algorithm that provides best solution for Steiner tree problem.

T–M Heuristic algorithm: One of the best known heuristic forf Steiner problem is T–M algorithm [6]. The algorithm operates in a manner that is similar to Prim's algorithm for MST. At each step a receiver is added to the tree. The receiver that is added has the shortest path between itself and the currently existing tree, just as the node that is added in the algorithm has the shortest path. The difference between the two procedures is that the path in the algorithm is a single link, allowing a straight forward search, while the path in the T–M heuristic algorithm may contain several links.

Modified T–M Heuristic algorithm: An important difference between the problem of multicasting and the conventional Steiner tree problem is that the cost of a link is not fixed but depends upon the maximum rate of the receivers that share the path. To solve this problem Maxemchuk [10] has proposed a heuristic called Modified T-M algorithm. In Modified T-M algorithm, we first form a Steiner tree with a high rate set of receivers and then successively add lower rate sets of

receivers to that tree with the same procedure as for the higher rate set.

3 M-MENTOR Algorithm

From the principles outlined in the previous section, we propose a modified MENTOR-II algorithm, called M-MENTOR algorithm, for IP networks with following features:

(1) All network members are within the same AS,

(2) The employ routing protocol should support multiple link weight such as M-OSPF and M-ISIS,

(3) The multicast traffic from different source share the same multicast tree. e.g. PIM-Bidir is deployed.

M-MENTOR starts with node clustering and select the backbone nodes in the same way as MENTOR and MENTOR-II.

Next, since all multicast transmitters share the same tree to distribute the data, instead of building Prim-Dijkstra tree, the backbone spanning tree is build based on Modified T-M algorithm. The algorithm forms a Steiner tree with a high rate set of receivers and then successively adds lower rate sets of receivers to the tree and, finally, adds nodes without receiver (zero rate receivers) to the tree. To guide the multicast traffic to flow only through the core spanning tree the multicast link weight should be set appropriately. One simple solution is to set the multicast link weight of all links in the spanning tree to 1 and that of the others to very large number.

After we obtained a Steiner tree, direct links are added and their unicast link weights are determined based on the unicast traffic in exactly the same way as in MENTOR-II.

The capacity a link on the tree is determined by the sum of its own unicast traffic, the overflow unicast traffic from other routes and the multicast traffic that flow through it.

4 Design Example

4.1 Requirements

An organization network composes of 6 backbone nodes shown in Figure 1. The network must support both unicast and multicast traffic. Table 1 and Table 2 show the multicast and unicast traffic demands between backbone nodes in the network. Table 3 shows the installation cost of link with 64 Kbps capacity. It is decided that the reasonable range of link utilization is determined by $\rho = 0.5$ and slack s = 0.4

S\R	N1	N2	N3	N4	N5	N6
N1	-	256	-	128	128	256
N2	256	-	-	128	128	256
N3	-	-	-	-	-	-
N4	128	128	-	-	128	128
N5	128	128	-	128	-	-
N6	256	256	-	128	128	-
					т	·/ TZ1

Table 1: Multicast traffic between backbone nodes

Unit: Kbps

Table 2: Uniicast traffic between backbone nodes

S\D	N1	N2	N3	N4	N5	N6						
N1		6120	1421	684	3472	5302						
N2	6120		10849	685	3629	4717						
N3	1421	10849		990	5863	4878						
N4	684	685	990		7267	5376						
N5	3742	3629	5863	7267		24747						
N6	5302	4717	4878	5376	24747							
					U	Jnit: bps						

Table 3: 64 Kbps Channel Link Installation Cost

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S	D	N1	N2	N3	N4	N5	N6					
NI	L	500	5905	8255	6915	6720	5740					
N2	2	5905	500	7955	8385	7960	7900					
N3	3	8255	7955	500	3185	2720	3985					
N4	1	6915	8385	3185	500	1160	1710					
Nź	5	6720	7960	2720	1160	500	1800					
Ne	5	5740	7900	3985	1710	1800	500					

4.2 Steiner Tree Construction

From Table 1, the backbone node can be classified into 3 groups based on capacity of the receiver, e.g. 256, 128 and 0 Kbps (no receiver), as follows.

G ₂₅₆	$= \{N1, N2, N6\}$
G ₁₂₈	$= \{N4, N5\}$
G_0	$= \{N6\}$

M-MENTOR first constructs a Steiner tree to interconnect members of G_{256} , i.e. N1, N2 and N6. At this stage, the installed links are N1-N2 and N1-N6 Next members of G_{128} , i.e. N4, N5, are connected to the tree with N4-N6 and N4-N5. Finally, member of G_0 , i.e. N6, is connected to the tree with N3-N5.

4.3 Multicast Traffic

In this study we assume that each transmitter open one different multicast session that cannot be mixed with other session.

Consider N1-N2 link, N2 may receive 2 of 256 Kbps multicast sessions and 2 of 128 Kbps multicast sessions simultaneously. Thus total multicast traffic on

N1-N2 direction is 768 Kbps. On the other hand, N2-N1, N1 may receive one of 256 Kbps multicast session.

Consider N6-N1 link, N1 may receive 1 of 256 Kbps multicast and 2 of 128 Kbps multicast sessions simultaneously. Thus total multicast traffic on N6-N1 direction is 512 Kbps. On the other hand, N1-N6, N6 may receive 2 of 256 Kbps multicast session simultaneously. Thus total multicast traffic on N1-N6 direction is 512 Kbps.

Consider N4-N6 link, N6 may receive 2 of 128 Kbps multicast sessions simultaneously. Thus total multicast traffic on N4-N6 direction is 256 Kbps. On the other hand, N6-N4, N4 may receive 3 of 128 Kbps multicast sessions simultaneously. Thus total multicast traffic on N1-N6 direction is 384 Kbps.

Consider N5-N4 link, N4 may receive 1 of 128 Kbps multicast session. On the other hand, N4-N5, N5 may receive 4 of 128 Kbps multicast sessions simultaneously. Thus total multicast traffic on N1-N6 direction is 512 Kbps.

According to Table 1 there is no multicast traffic on N3-N5 link.

4.4 Overall Traffic and Design Results

Using ISP algorithm as in MENTOR-II, 2 new direct links N2-N5 and N5-N6 are installed for unicast traffic. Figure 1 and Table 4 show the final design result obtained by M-MENTOR. In Table 4, "weight" is the unicast weight of the link obtained by ISP algorithm; "load" is the total traffic on each direction of the link; "# Chan" is the number of 64 Kbps channels of the link. Traffic load on each direction may compose of unicast and multicast traffic. If available, the number represented the amount of multicast traffic is underlined.

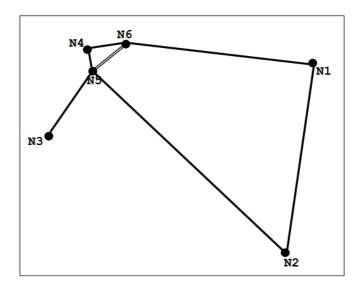


Figure 1 : Network Design Result

			8	
Link	s_i, d_i	weight	Load	#
		(unicast)	(bps)	Chan
1	N1,N2	5,905	6,120+ <u>768,000</u>	25
	N2,N1		(6,120+ <u>256,000</u>)	
2	N1,N6	5,740	10,879+ <u>512,000</u>	17
	N6,N1		(10,879+ <u>512,000</u>)	
3	N4,N6	1,710	6,060+ <u>256,000</u>	13
	N6,N4		(6,060+ <u>384,000</u>)	
4	N4,N5	1,160	8,942+ <u>512,000</u>	17
	N5,N4		(8,942+ <u>128,000</u>)	
5	N5,N3	2,720	24,001	1
	N3,N5		(24,001)	
6	N2,N5	8,774	19,880	1
	N5,N2		(19,880)	
7	N5,N6	2,869	39,235	2
	N6,N5		(39,235)	
	•	•	·	•

Table 4: M-MENTOR Design Results

5 Performance Evaluation

In order to evaluate the efficiency of network design calculated by M-MENTOR algorithm, we analyze the performances of a number synthesized network and in term of installing cost.

5.1 Requirement Generation

To evaluate the efficiency of M-MENTOR algorithms for different sizes, three design requirements of 10, 30 and 50 nodes are generated. A design requirement composes of node distribution and the associated traffic demand matrix. For each requirement, a design tool called DELITE [5] is used to synthesize nodes location distribution which is obtained by randomly varying SEED parameter of DELITE. All node distributions have average node distances of around 800 kilometers and maximum node distance of around 1600 – 1900 kilometers. The unicast traffic demand for each requirement is also generated by DELITE with following assumption:

1) All nodes have the same total unicast traffic in and total unicast traffic out. In this study, the total

unicast traffic in and out of 256, 512, 1024, 2048, 3072 and 4096 Kbps are considered

2) The unicast traffic between a pair of node is inverse proportional to the distance between them.

In this study we assume that a link interconnecting a pair of nodes could be installed with multiple communication channel of bandwidth 64 kbps. The link installation cost is linearly proportional to the distance between nodes and the number of 64 kbps channel on the link. For the 10, 20 and 30 nodes requirements tables 5, 6 and 7, respectively, show the link installation cost with single 64 kbps channel generated by DELITE.

For a given amount of unicast traffic demand, various levels of multicast traffic demands are considered. For simplicity, in this study, we use the notation $M \ge 512 + N \ge 256$ which denotes that, at the same time, M + N nodes of alls are generating different multicast session where M of them are 512 kbps session and the other N are 256 kbps session.

5.2 Network Design Results

Tables 8 - 22 present the installation cost for 10, 30 and 50 nodes networks obtained by M-MENTOR with $\rho = 0.4$, slack s = 0.2, 0 .4 and 0.6, and the total unicast traffic in/out per node of 256, 512, 1024, 2048 and 4096 kbps. Since there is no known algorithm for design of mesh IP network with mixed traffic before, for comparison, we have modified MENTOR-II to support both unicast and multicast traffic. What we have modified is that all multicast traffics are forced to flow through the backbone spanning tree of the MENTOR-II network. And thus the link capacities of the spanning tree are determined by the summation of the unicast and multicast traffic flow through them. The design results obtained by modified MENTOR-II with $\alpha = 0$ (i.e. Minimum Spanning Tree: MST), 0.5 and 1 (i.e. Shortest Path Tree: SPT) are also presented in the tables in order to compare with of network obtained by M-MENTOR that have the same slack

S\D	N1	N2	N3	N4	N5	N6	N7	N8	N9	N10
N1	500	1174	830	2340	2860	3504	1236	3122	1302	1718
N2	1174	500	994	1870	1182	2514	3592	3506	3348	1820
N3	830	994	500	2328	1502	2832	2006	2434	3546	2184
N4	2340	1870	2328	500	1450	2760	2912	1662	1672	2330
N5	2860	1182	1502	1450	500	2560	2690	2228	2124	1998
N6	3504	2514	2832	2760	2560	500	2716	1770	1660	1470
N7	1236	3592	2006	2912	2690	2716	500	2442	1360	1386
N8	3122	3506	2434	1662	2228	1770	2442	500	1956	1020
N9	1302	3348	3546	1672	2124	1660	1360	1956	500	1380
N10	1718	1820	2184	2330	1998	1470	1386	1020	1380	500

Table 5: 64 Kbps Channel Link Installation Cost for 10 Node Requirement

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Table

N30	1314	2370	2808	060	2140	2560	2780	3272	1890	2898	3430	3328	2250	2210	2698	2138	1504	804	1778	2094	2252	2924	428	3180	1200	056	244	1736	1248	ŰŞ
N29 1	1578 1	2014 2	736 2	2634 3090	2446 2	2422 2	2872 2	1508 3	2148 1	348 2	3626 3	1646 3	744 2	2364 2	1982 2	014 2	492 1	1804 8	1892 1	514 2	2274 2	3516 2	3464 3428	660 3	2726 1	134 3	1384 1244	1434 1	500	-
N28 P	940 1	2364 2	324 2	492 2	2118 2	2400 2	1604 2	1120 1	3276 2	000 3		1530 1	2108 1744	2124 2	2806 1	484 2	234 1	2012 1	798 1	448 1	3646 2	3840 3		2392 6	1862 2	626 1	2176 1	500 1		N30 1314 2370 2808 3000 2140 2560 2780 3272 1890 2808 3430 3328 2250 2210 2698 2138 1504 804 1778 2004 2252 2624 3428 3180 1200 3056 1244 1736 1248
N27 P		3702 2	2724 3324 2736	2246 2492	2422 2		878 1	3424 1		458 4	1784 1206	2260 1.	412 2		2920 2	586 2	304 2	1138 2		854 1	4124 3	2158 3	482 1	284 2	1994 1	2588 1626 1134 3056	500 2	2176 5	384 1	744 1
N26 P	798 2	1294 3			1746 2	2102 1942		2718 3	716 4	500 1.		1502 2	702 2	960 3	1926 2	512 2	102 2	1608 1	080 1	592 3	2872 4	580 2	356 2	892 2	2390 1	500 2			134 1	12
N25 P	006 2	2604 1	024 8	274 7	3592 1	492 2	784 1		518 1	516 1	946 8		836 2	632 2	2020 1	068 1	336 2	2180 1	664 2	624 2	726 2		210 2	714 1	20		994 2	862 1	726 1	2000
N24 P	1148 2006 2798 2438	1522 2	488 3	476 2	1406 3	2086 2492	792 1	1464 2788	2462 1826 2922 1226 1074 2518 1716 4280	964 2	2138 2658 1974 1946 892	2366 1146	286 1	058 1	2196 2	070 3	656 3	1938 2	2546 1664 2080 1368	500 1306 2054 2044 1222 2624 2592 3854 1448 1514	1904 2726	2092 1576 2040	814 2	500 1714 1892 2284	1714	2356 1892 2390	3854 4124 2158 2482 2284 1994 2588	2392 1862 1626	560 2	100
N23 I		1430 1	940 1	680 2	2698 1		194 1	594 1	226 1	444	658 1		220 1	840 1	480 2	144 2	558 1	2558 1	806 2	044 1	1576 1	092 1	500 1		2210 1	356 1	482 2	1726 2	:464	000
N22	904	2066 1	666 1	912 1	1798 2	628 1	084 1	1712	922 1	154 2	138 2	1714 1650	156 1	562 1	2242 1480	792 2	014 2	802 2	2532 8	054 2	2588 1	500 2	092	576 1	2040 2	580 2	158 2	3840 1	516 3	200
N21	3112 1822 1904 998	2372 2	906 1	020 2	832 1	2222 1362 1628 1438	052 1	854 1	826 2	148 3	942 2	1320 1	642 2	8626 1	2744 2	634 2	728 3	1148	654 2	306 2	500		576 2	904 1	2726 2		1124 2	3646 3	274 3	5
1 22 2	3112 1	4054 2	304 2	594 1	1968	222 1	8160 1	2962 1854	3462 1	122 1		1006	3576 1	3902 2	2304 2	2154 1	286 1	1910 1	2558	500	1306	2054 2588	2044 1	222 1	2624 2	292 2	854 2	1448	514 2	
N19		2576 4	2274 1	l 672 1	1868 1		3208 3	2156 2		2988	2630 1	1116 1	1798 3	3450 3	1474	2804 2	2472 2	2684 1	8		654 1	2532	806	2546 1	1664	2080	1368 3	798 1	1892	011
N18	3326 1878	3000	2574 2154 1456 2638 2274 1304 2906 1666 1940 1488 3024 818	614 1558 2156 2368 2378 1672 1594 1020 2912 1680 2476 2274 728	2342	1824 1030 1660 3034 2700	2346	2488	2618 1176 3706 1958	<i>6</i> 78 1460 1452 2494 2842 2942 3652 2988 2122 1148 3154 2444 964 2516 1500 1458 4000 3348	1886 3170 1284 2630 1148	662]	1408 1506 1472 1732 3662 1798 3576 1642 2156 1220 1286 1836 2702 2412	980 1554 2640 3450 3902 2626 1562 1840 1058 1632 2960 3212	2992 1474	500 1788 2922 2804 2154 1634 2792 2144 2070 3068 1512 2586 2484 2014	2718 2	500	2684	006 3576 3902 2304 2154 2286 1910 2558	1148	802	650 1220 1840 1480 2144 2558 2558 806 2044 1576 2092 500 1814 2210 2356 2482 1726	2196 2070 1656 1938 2546 1222 1904 1576 1814	2180 1664	502 2702 2960 1926 1512 2102 1608 2080 2592 2872	1138 1	2012	1804	, VO
μŢ			1456	2368	2254	1660	1348	912	1176	2942	3170		1732	1554	1120	1788 3	500	2718	2472	2286		3014	2558	1656	3336	2102	2304	2234	1492	202
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24 23 24 23 29 29	4 2 2	1270 233 4502 449	1 22	30	1770 223	1 1	02 11	12	50	7	20 13 14 14	50	50	2472 280 1910 228	11	54 25	76 20	76 19(20 110	2390 1292 2356 580	34	26 13	: ÷	1 4	02 31	36 11	3070 2800 3168 1616 5336 4746 5536 835	1 12 1 12	50 10	74 17	1 1 2 2 2 2	10 20	2438 251	7 7 7 7 7 7 7 7 7 7 7 7 7 7 7 7 7 7 7	20 23	99 3 5 12	9 17 9 00	62	62 28
M1 M 500 11 1174 50	230 99	2348 18	2560 27	2716 26	2442 17	1320 10	11 12	70 16	34 17	11 12	32 13	22 11	112 29	2624 24 2558 19	06 63	20	92 15	14 15	14 22	190 12	19	176 16 34 13	7 7 7	1 1	96 27	19 29	2 X 2 X	2660 20	22 26	04 22	9 Q	2128 16	2254 24	2154 22 262 15	2528 18	2344 25	2182 17	1962 14	16 25
N 20 N		114 53 Mit 23		5 5	21 ST	3110	201 1032	112 10	112	7	112 112 112 112	11 12	118 27	2019 26 2020 25	3721 13	22	22	124 12	125 17	3126 23	2	122 21 1070 14	10	1 G 1 E	152 31	5	1024 30 Wee 13	B12 (50	M57 1022	158 17	139 11 11 11	3144 24	3142 22	N45 21	3145	B146 23	24 T42 24	1149	20 53
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Table 8: Installation cost of 10 node network with tota	1
unicast traffic in/out of 256 Kbps	

Multicast	slack	M-MENTOR	MENTOR-II										
Traffic	SIACK	WI-WIENTOK	$\alpha = 0$	$\alpha = 0.5$	$\alpha = 1$								
2x512+	0.2	10430	11400	13616	15140								
6x256	0.4	11684	12160	14226	15772								
	0.6	12248	12864	15905	16834								
4x512+	0.2	12862	13374	15457	16957								
4x25	0.4	14438	15976	16885	17206								
	0.6	16825	18157	19215	19848								
6x512+	0.2	14045	14365	16125	17426								
2x256	0.4	16526	18853	19596	19632								
	0.6	19986	20268	20904	21159								

Table 9: Installation cost of 10 node network with total unicast traffic in/out of 512 Kbps

Multicast	slack	slack M-MENTOR	MENTOR-II		
Traffic	STACK	WI-WIENTOK	$\alpha = 0$	$\alpha = 0.5$	$\alpha = 1$
2x512+	0.2	12216	12537	13892	15524
6x256	0.4	14560	14763	16162	17873
	0.6	15656	16930	17483	19201
4x512+	0.2	13463	13834	15314	17027
4x25	0.4	16862	17164	18855	19480
	0.6	17289	18583	19238	20772
6x512+	0.2	15320	15376	16993	17738
2x256	0.4	16756	17860	18423	19912
	0.6	20952	21027	21829	22114

Table 10: Installation cost of 10 node network with totalunicast traffic in/out of 1024 Kbps

Multicast	slack	M-MENTOR	MENTOR-II		
Traffic	SIACK	WI-WIENTOK	$\alpha = 0$	$\alpha = 0.5$	$\alpha = 1$
2x512+	0.2	14728	14920	16378	17902
6x256	0.4	16803	17209	18839	19328
	0.6	17264	18678	19238	20728
4x512+	0.2	16172	17427	18537	19392
4x25	0.4	17489	18874	19839	20795
	0.6	18638	18928	19299	20945
6x512+	0.2	16284	17538	18783	19458
2x256	0.4	17620	18947	19944	20722
	0.6	20114	21264	21326	22928

Table 11: Installation cost of 10 node network with total

 unicast traffic in/out of 2048 Kbps

Multicast	slack M-MENTOR	MENTOR-II			
Traffic	SIACK	WI-WIENTOK	$\alpha = 0$	$\alpha = 0.5$	$\alpha = 1$
2x512+	0.2	16138	16483	17672	18628
6x256	0.4	17433	18820	19930	20920
	0.6	19894	20281	21218	22219
4x512+	0.2	17829	18978	19839	20829
4x25	0.4	18945	19281	20281	21114
	0.6	20575	20839	22839	23438
6x512+	0.2	19118	19271	20921	21930
2x256	0.4	20439	20629	21218	22103
	0.6	21842	21920	22618	23543

Table 12: Installation cost of 10 node network with total
unicast traffic in/out of 4096 Kbps

Multicast	slack	M-MENTOR	MENTOR-II		
Traffic	SIACK	WI-WIENTOK	$\alpha = 0$	$\alpha = 0.5$	$\alpha = 1$
2x512+	0.2	20748	19849	22955	24184
6x256	0.4	20022	20294	23329	24499
	0.6	21493	21543	22849	25893
4x512+	0.2	22431	23452	25198	25384
4x25	0.4	23738	24873	26493	29594
	0.6	27839	29918	29899	30883
6x512+	0.2	23855	24890	25783	26739
2x256	0.4	25937	26182	28990	30922
	0.6	29211	31483	32352	33274

Table 13: Installation cost of 30 node network with total unicast traffic in/out of 256 Kbps

Multicast	slack M-MENTOR	M MENTOP	MENTOR-II		
Traffic	SIACK	WI-WENTOK	$\alpha = 0$	$\alpha = 0.5$	$\alpha = 1$
6x512+	0.2	34586	36270	44250	47410
18x256	0.4	35632	36459	44567	47974
	0.6	38473	38943	45362	48348
12x512+	0.2	36958	37845	45365	48594
12x25	0.4	38647	40432	47948	49320
	0.6	40865	43854	51844	52043
18x512+	0.2	40538	40538	48357	50384
6x256	0.4	41847	42843	51335	52367
	0.6	43562	46484	52847	53475

Table 14: Installation cost of 30 node network with totalunicast traffic in/out of 512 Kbps

Multicast	slack	M-MENTOR	MENTOR-II		
Traffic	SIGCK	WI-WIENTOK	$\alpha = 0$	$\alpha = 0.5$	$\alpha = 1$
6x512+	0.2	40372	41746	45374	48438
18x256	0.4	42847	43728	47384	50324
	0.6	47487	48739	50387	51472
12x512+	0.2	46378	48376	49578	51362
12x25	0.4	49384	50367	52748	53892
	0.6	52832	54633	54218	54362
18x512+	0.2	52738	52738	53127	53728
6x256	0.4	53627	53627	54372	55382
	0.6	53903	54637	55283	56372

Table 15: Installation cost of 30 node network with totalunicast traffic in/out of 1024 Kbps

Multicast	slack	ick M-MENTOR	MENTOR-II		
Traffic	SIACK	WI-WIENTOK	$\alpha = 0$	$\alpha = 0.5$	$\alpha = 1$
6x512+	0.2	48374	49302	49302	51894
18x256	0.4	49374	49374	50758	54378
	0.6	50489	51394	52948	55938
12x512+	0.2	49803	50549	51948	55281
12x25	0.4	51038	52849	53849	58394
	0.6	55847	57384	59847	61346
18x512+	0.2	52748	53472	54637	60483
6x256	0.4	55493	55493	57439	62849
	0.6	49083	61849	63584	64839

Multicast	slack	M-MENTOR	MENTOR-II		
Traffic	STACK	M-MENIOK	$\alpha = 0$	$\alpha = 0.5$	$\alpha = 1$
6x512+	0.2	50473	50473	51847	55374
18x256	0.4	50973	51849	53948	57483
	0.6	51756	52673	54890	58494
12x512+	0.2	50438	51874	59039	57849
12x25	0.4	53746	55849	60487	60382
	0.6	57684	58393	62847	63928
18x512+	0.2	57467	58948	59860	60485
6x256	0.4	59485	60494	62748	63859
	0.6	61857	62847	65489	66854

Table 16: Installation cost of 30 node network with total unicast traffic in/out of 2048 Kbps

Table 17: Installation cost of 30 node network with total unicast traffic in/out of 4096 Kbps

Multicast	alaala	slack M-MENTOR	MENTOR-II		
Traffic	STACK	WI-WIENTOK	$\alpha = 0$	$\alpha = 0.5$	$\alpha = 1$
6x512+	0.2	53467	53467	58475	60284
18x256	0.4	55768	56372	59485	62837
	0.6	60574	61738	63849	64738
12x512+	0.2	62547	62547	62547	63748
12x25	0.4	62768	63546	64758	66478
	0.6	63758	64564	67483	68476
18x512+	0.2	62786	63758	65374	67463
6x256	0.4	64768	65867	67645	69478
	0.6	66970	67869	69384	71647

Table 18: Installation cost of 50 node network with total unicast traffic in/out of 256 Kbps

Multicast	slack	M-MENTOR	MENTOR-II		
Traffic	STACK	WI-WIENTOK	$\alpha = 0$	$\alpha = 0.5$	$\alpha = 1$
10x512+	0.2	38134	39690	43890	59650
30x256	0.4	38845	40495	44192	60581
	0.6	39520	41345	45648	61378
20x512+	0.2	39148	40982	44357	60964
20x25	0.4	39850	41385	45856	61851
	0.6	41183	42724	46364	62774
30x512+	0.2	40643	41638	46032	62157
10x256	0.4	41749	42581	46790	62814
	0.6	42752	43485	47371	63825

Table 19: Installation cost of 50 node network with total unicast traffic in/out of 512 Kbps

Multicast	slack	M-MENTOR	MENTOR-II		
Traffic	STACK	WI-WIENTOK	$\alpha = 0$	$\alpha = 0.5$	$\alpha = 1$
10x512+	0.2	39231	40712	44061	60824
30x256	0.4	40574	41664	44945	61367
	0.6	41259	42305	46123	62231
20x512+	0.2	40532	41563	44926	61247
20x25	0.4	41576	41903	45945	62394
	0.6	42183	42945	47190	63283
30x512+	0.2	42227	42594	46583	62954
10x256	0.4	42879	43186	47243	63436
	0.6	43472	43975	48521	64221

Table 20: Installation cost of 50 node network with total unicast traffic in/out of 1024 Kbps

Multicast	slack	M-MENTOR	MENTOR-II		
Traffic			$\alpha = 0$	$\alpha = 0.5$	$\alpha = 1$
10x512+	0.2	40365	41475	44724	61428
30x256	0.4	41634	42193	45638	62248
	0.6	42254	42978	46883	63482
20x512+	0.2	41749	42532	45931	62185
20x25	0.4	42364	42954	46365	63854
	0.6	42698	43367	47811	64284
30x512+	0.2	42786	43276	47365	63265
10x256	0.4	43458	44675	47956	64462
	0.6	44284	45834	49257	65398

Table 21: Installation cost of 50 node network with total unicast traffic in/out of 2048 Kbps

Multicast	slack	M-MENTOR	MENTOR-II		
Traffic			$\alpha = 0$	$\alpha = 0.5$	$\alpha = 1$
10x512+	0.2	41685	42256	42256	62456
30x256	0.4	42490	43678	43898	63648
	0.6	43643	44394	45274	64267
20x512+	0.2	43574	43854	46426	63389
20x25	0.4	43685	44253	46852	64286
	0.6	43894	44845	48258	65378
30x512+	0.2	43584	43857	47899	63992
10x256	0.4	44367	45684	48538	65170
	0.6	45336	46936	49703	66357

Table 22: Installation cost of 30 node network with total unicast traffic in/out of 4096 Kbps

Multicast	slack	M-MENTOR	MENTOR-II		
Traffic	SIDEK		$\alpha = 0$	$\alpha = 0.5$	$\alpha = 1$
10x512+	0.2	43056	43224	42804	63109
30x256	0.4	43857	44217	44730	64743
	0.6	44721	45134	45997	65807
20x512+	0.2	44652	44906	47255	64423
20x25	0.4	44978	45354	47658	65284
	0.6	45143	45852	48957	65920
30x512+	0.2	44365	44707	48865	64889
10x256	0.4	44795	45290	49573	65902
	0.6	45644	46456	50965	67348

It can be observed from tables 8 - 22 that:

- 1) Given a value of slack, almost all cases M-MENTOR gives the minimum installation cost. On the other hand, MENTOR-II, the installation cost increase as α increase.
- 2) Given a traffic demand, the installation cost tends to increase as slack increase.
- 3) As traffic demands increase, either unicast or multicast, the installation cost increase.

6 Conclusion

This study proposes an IP network design algorithm called M-MENTOR that support both unicast and multicast traffic simultanously. However, since multicast traffic model could be employed in many situations and could be managed by various kinds of multicast routing protocols, this paper consider only the IP networks with following features: (1) network of within the same Autonomous System (AS), (2) routing protocol support multiple weight on each link, (3) the multicast traffic from different sources share the same multicast tree. M-MENTOR is a modified version of MENTOR-II that uses Modified T-M algorithm, rather than Prim-Dijkstra Algorithm, to construct backbone spanning tree. An example of 6 backbone nodes network design is given.

The efficiency of M-MENTOR is evaluated in term of network installation cost. The installation cost of 10, 30 and 50 backbone nodes networks designed by M-MENTOR are calculated and compared with that of MENTOR-II with various design parameters and various conditions of mixed traffics. It is shown that, all most all cases, M-MENTOR networks give lowest installation cost. For the case of MENTOR-II, networks with lower α tend to give better performance.

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