HMI and Control Loops over Serial Line Communication Channel with MODBUS protocol for Combinatorial Weigher at Potato Chips Production Line

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Abstract: - Potato chips are product with irregular forms with complex flow dynamics. Packing with narrow tolerance is done by combinatorial weighing machine. For controlling this complex machine it is of great importance to have well designed human-machine interface (HMI). To minimize human intervention on machine operation of great importance is quality of its control algorithm. To gain an insight into the process of vibratory conveying and combinatorial weighing of potato chips experimental on-line setup was created. Setup contains refurbished combinatorial weigher and personal computer (PC) with network control structure. First, a short introduction to the chain of supplying, weighing and packing at potato chips production line is given. It is described briefly principle of combinatorial weighing. Next, experimental setup was described, from constructive parts, to hardware and software implementation. RS485 serial communication line with MODBUS protocol was used for networking master and three slave control units. Application software and implemented functions for making flexible development system are presented. HMI and MODBUS protocol were implemented with Visual Basic. Finally, tests results are given and show flexibility of setup in implementing real-time closed-loop control over MODBUS communication channel for combinatorial weigher at potato chips production line.

Key-Words: - Combinatorial Weigher, Potato Chips, Vibration Feeder, HMI, MODBUS, Visual Basic

1. Introduction
Most of food manufacturing industries as last part of production line have some sort of packing machine. Quality demands on producer side now days are high. One of mandatory requests is real pack weigh and one marked on package and "at-least" rule applies in this case. On the other side, "over-weighted" packages are in disfavor viewed from producer side. Production line is continuous chain of machines, from preparation of raw materials to handling of already packed final product. Throughput of each segment of the chain must be compatible with requested capacity. Weighing final product for packing with necessary speed, while fulfilling "at-least" rule and minimizing "over-weight" part, imposes specific measuring equipment. This part of process is even harder if size and weight of product particles varies in broad range and with complexity of flow mechanism of products with irregular forms. Member of this category are Potato Chips range of products. Solution for weighing this category of products is Automatic Combination Weigher [1], or Multihead Weigher [2] [3]. Core of combinatorial weighing is acquiring weigh data from multiple weighing heads, calculating sums of all possible combinations and selecting the best one, i.e. combination that gives sum weigh closest to "at-least" weigh from upper side [4] [5]. Second designation, multihead weigher, must imply combinatorial process of weigh data to be part of this grup. Motivation for this work was complexity of dynamic of this simply stated principle of operation and lack of literature that treats whole process from product supply to weighted pack product. Aim was to create a platform for gathering data from real process, for modelling and simulating [6], and for testing different control strategies that will improve accuracy and speed, reduce downtime and giveaway and in the same time having all the functionality of modern Human-Machine Interface (HMI) [7]. Since control loops will go through communication channel it will be a good platform for testing industrial communication and real-time control over network with Master-Slave topology. Selected protocol was MODBUS over Serial Line [8] [9]. As a physical layer 2-Wire RS485 was adopted, as
industry widely used bidirectional, balanced transmission line standard [10].

2. Supplying, Weighing and Packing Chain

Potato chips production line uses potatoes as incoming raw material. After preparation phase, potatoes get sliced up. Slices enter the frying oven where they get fried and dried of excessive fry fat. Next stage is visual inspection and removal of out of tolerance fryed and other defective slices. Flavoring is done in next step and again followed with visual inspection for proper level of flavoring and if necessary removal of over flavored. At this stage product is ready for consumption and for quality of product it important to reduce the time of contact product- free air. Here starts the chain of Supplying-Weighing-Packing. Continuous flow of product from previous stages is delivered to supplying conveyor, Fig.1, for transporting it to the entrance of combinatorial weigher. Weighing, finding the best final weight and making it ready for packing machine is duty of combinatorial weighing machine. Final weight that is waiting will be released to packing machine on its request.

2.1 Supply Conveyor

Transporting product from near floor level to near ceiling level to supply position of combinatorial weigher is done with conveyor of bucket type. Its state of operation is controlled by the level of product in receiver side, conical part marked by CF (Circular Vibratory Feeder) in Fig.1. Conveyor is driven by Induction AC motor and usually is controlled by simple On/Off algorithm. For improved performance and reduction of wear-off motor was supplied by Variable Frequency Drive (VFD) [11]. VFD is equipped with local control panel and with MODBUS module. It is one of slaves in the network.

2.2 Combinatorial Weigher

Weigher is composed of three functional groups, Fig.1 . First, upper group, is product supply side. Main part of this group is central conical body that's vibrating surface of central vibratory feeder (CF). In middle there are \( N_h \) measuring heads. Each head has one horizontal radial vibratory feeder (RF), one holding hopper (HH) with controllable door at its bottom, one weighing hopper (WH) with two independently controllable doors (left and right), and two memory hoppers (MH). Third function group is exit unit with collecting chute and discharge hopper (DH) as its parts. DH has controllable door for discharging product on request from packing machine. CF serves as a distributor of incoming product to \( N_f \) heads. Rate of product flow is controlled by intensity of vibrations. Product that leaves CF falls to corresponding radial vibratory feeders surface. Perimeter of CF is divided in \( N_h \) equal sectors. Product on surface of RF is transported and supplied systematically [12] [13] to HH. When weigh is free for new measurement, content of HH as a batch is released to WH and weight data are acquired. If one of MHs is empty, corresponding door is opened and product goes to memory hopper. Weight data from \( 2*N_h \) MHs serve as input to combinatorial process [4]. In this process are calculated sums of all possible combinations \( \binom{N_h}{k} \) for \( k=1 \) to \( h \). Results are compared with set weight and with maximal tolerable weight. From values in this range is selected one closest to the set weight. Selected value is sum of k weights from selected MHs. If DH is empty selected MHs are opened and their product is collected to DH. Now exit unit is ready for supplying a batch of product to packing machine. Packing machine is informed and unit waits for request signal. When exit unit receives request from packing machine it will open doors of DH from where product will go to prepared pack. Upon release of product exit unit will be ready for new batch.
3. Experimental On-Line Setup

For creation of experimental setup a second handed automatic combination weigher was used. Weigher was an old model, K723A [14], of Anritsu [1]. It is an 8-head weighing machine. Declared speed was 60 to 120 packs/minute. Weighing range is 50g to 2000g and accuracy $\pm 1$ to 1.5g. All specifications are dependent on product and installation conditions.

Mechanical, pneumatic and electromagnetic vibratory feeders were tested with positive result. Electronics was defective and it was found that almost all EPROMs had experienced aging caused data corruption. Analog and digital I/O units (CH1~CH4, CH5~CH8, A/D, CONT, I/O, F.DRIVE) were functional and used as such. ALU unit that was responsible for combinatorial calculations was not used. Z80CPU unit is dual microprocessor board with 2 Zilog Z80H-CPU microprocessors operating at 8MHz clock [15]. They have possibility of using shared-memory block, which is not used. Two microprocessors now communicate only over network. One functions as a Slave1 and other as a Slave2. Keyboard and display (KEY&DISP) was not used too.

Functions of combinatorial calculations and HMI are done with a Personal Computer (PC). PC is Master of the network. It is an PIV PC at 1GHz, with 512MB RAM and 20GB HDD.

3.1 Network Topology

Physical layer of network is implemented as "2-Wire" electrical interface in accordance with EIA/TIA-485 standard [10]. Only one driver can be active at any time. Network topology with all actors is given in Fig.2.

It is a multipoint serial line bus and requirements of [8] are fulfilled. Connection of Line Termination (LT) impedances (resistor 150Ω/0.5W) are placed at both ends. Line polarization resistors PU (Pull-Up) and PD (Pull-Down) with value 650Ω were used. Three passive taps with two IDv (Distributor type Derivation Interface) were used. First tap is used for connection of Slave1 (CPU1) and Slave2 (CPU2). Signal ground connection to protective ground and one of LT are placed at this tap. One IDv of second tap is used for connection of Slave1 (CPU1) and Slave2 (CPU2) to bus. Third tap is connection point of Master to bus. Second LT is placed at this tap. Unconnected second IDv of last two taps may be used for diagnostic monitoring activities on bus.

Trunk cable lengths were 10m and 30m and are not critical for used Baud Rate of 19200bps.

3.2 Protocol and Data Model

Selected Protocol was MODBUS over serial line in RTU transmission mode. If compared to OSI model with 7 layers, this protocol takes place at level 2 (Data Link) and level 7 (Application) [8]. Protocol provides client/server communication. Protocol Data Unit (PDU) is defined by protocol specification. Additional fields are introduced to create Application Data Unit (ADU) and this part is bus and network dependent.

PDU has Function code field and Data field. Over serial line communication appends preceding Address field and ending CRC field. ADU in this case is referred as "MODBUS frame over Serial Line" [8].

Frame for RTU mode is shown in Fig.3. Address field has one byte. Valid slave numbers are 1-247. One byte Function code field defines type of request from client. There are three types of requests: Public, User-Defined and Reserved [9] [16]. Public

<table>
<thead>
<tr>
<th>Field</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>Start</td>
<td>$\geq$ 3.5 char</td>
</tr>
<tr>
<td>Address</td>
<td>8 bits</td>
</tr>
<tr>
<td>Function</td>
<td>8 bits</td>
</tr>
<tr>
<td>Data</td>
<td>Byte 1 x 8 bits</td>
</tr>
<tr>
<td></td>
<td>Byte 2 x 8 bits</td>
</tr>
<tr>
<td></td>
<td>...</td>
</tr>
<tr>
<td></td>
<td>Byte N x 8 bits</td>
</tr>
<tr>
<td>CRC Check</td>
<td>16 bits</td>
</tr>
<tr>
<td>End</td>
<td>$\geq$ 3.5 char</td>
</tr>
</tbody>
</table>

Fig. 3 RTU message frame
functions are well defined, unique and publically documented. Function defines structure and number of bytes in Data field and it is variable. CRC field is error checking field and is result of "Redundancy Checking" calculations. It is 2 bytes wide. Silence interval $\geq 3.5$ character times serves as frame separator. Inter-frame allowed silence is $\leq 1.5$ character times. This time intervals are marked $t_{3.5}$ and $t_{1.5}$. In our case $t_{3.5} = 1/19200*11*3.5 = 2\text{ms}$ and $t_{1.5} = 1/19200*11*1.5 = 0.859\text{ms}$.

Data abstraction in MODBUS is done by defining four primary tables with defined Object type and Type of access, Table 1.

<table>
<thead>
<tr>
<th>Primary table</th>
<th>Object type</th>
<th>Type of access</th>
</tr>
</thead>
<tbody>
<tr>
<td>Discrete Input</td>
<td>Single bit</td>
<td>Read-Only</td>
</tr>
<tr>
<td>Coils</td>
<td>Single bit</td>
<td>Read/Write</td>
</tr>
<tr>
<td>Input Registers</td>
<td>16-bit word</td>
<td>Read-Only</td>
</tr>
<tr>
<td>Holding Registers</td>
<td>16-bit word</td>
<td>Read/Write</td>
</tr>
</tbody>
</table>

Table 1. Four primary tables of MODBUS.

Address for data in each table is from 0 to 65536. Mapping from data address of MODBUS data model to device application is device specific and can be freely selected. Overlapping between tables is possible, giving us e possibility of accessing same data in different formats.

### 3.3 Slave1 and Slave2 Application Software

Microprocessor for both slaves is Z80H-CPU, member of 8-bit family of microprocessors initially produces by Zilog. They are operating at 8MHz clock. Hardware system is has 8kB RAM/32kB EPROM and possibility of using another 8kB RAM block as a shared memory for both CPU systems.

Applications for slaves were developed in Assembly Language. Structure of software is based in stateflow as very convenient model for developing real-time event-driven software [17]. Model followed during development was subset that resembles modelling with Stateflow toolbox in Matlab/Simulink environment [18].

EPROM was programmed with following blocks of software:

- EPROM/RAM CRC Tests
- Hardware Tests
- Initialization
- Watchdog (software part)
- MODBUS RTU protocol
- MODBUS Function implementation
  - 01 Read Coils
  - 02 Read Discrete Input
  - 03 Read Holding Registers
  - 04 Read Input Registers
  - 05 Write Single Coil
  - 06 Write Single Register
  - 15 Write Multiple Coils
  - 16 Write Multiple Registers
  - 23 Read/Write Multiple Registers

Application part
- CPU1: - Main Weighing
  - Weight Measurement
  - Weigh Calibrations
  - Read Temperatures
  - Low-Level Tests
- CPU2: - Circular Feeder Control
  - Radial Feeder Control
  - Hopper Control
  - Conveyor Control
  - Handshaking with Packer
  - Low-Level Tests
  - Firmware Development

Firmware Development is used during development of new software for slaves. During this phase software is loaded in RAM through MODBUS channel. Functions from RAM can be called and test for functionality. Content of RAM is protected with battery, but for safety reason mature versions are reprogrammed in EPROM.

Following the RESET, CPU tests EPROM and RAM content with same CRC algorithm that is used for MODBUS protocol. Hardware tests of peripherals are done next. If tests are passed with positive result Initialization phase is executed and Watchdog is started. CPU will listen on bus for MODBUS requests and will respond to those with correct address.

Watchdog is implemented in hardware and waits for periodical attention from CPU. It is physically under control of CPU2. In software part two levels of watchdog were implemented. First level is for watching for failures in handshaking between CPU1 and CPU2 through shared memory. Currently this is only usage of shared memory resources. Second level is watching handshaking of CPU1 / Master / CPU2 chain over MODBUS communication channel. Used level is user selectable.

MODBUS RTU protocol and Function execution is part of interrupt service routines of serial port and periodic timer. Events generated from serial port and timer guide transitions of stateflow. After receiving error-free request it is executed and response is sent to Master.

Application part runs in Main body of program. Transitions of main program may be generated from events and conditions of corresponding functional unit that it controls. Higher levels of control may
influence transition process by data written on specific registers.

CPU1 has 9 processes working in time-shared bases. Global process follows requests from Master for operation conditions:

- Stop
- Start
- Pause
- Continue
- Unit Tests
- Maintenance

Each of 8 heads has its own process and performs operation in accordance with Global process.

CPU2 has 12 parallel processes. 9 are same as those of CPU1. Other 3 are for controlling conveyor, circular feeder and handshaking with packer.

Global process of CPU1 and CPU2 are copies of Global process in Master. Master must receive acknowledgement from both slaves to proceed in new Global state, otherwise error is generated.

Under normal operation interaction between processes in different level is done with specific data in one of declared data spaces. For testing purpose, by using low-level functions, master may read or change any data in Slave memory space. This gives the opportunity of any type of control from Master by using small set of primitive functions. In lowest level Slaves serve only as raw MODBUS RTU Input/Output unit with no additional processing. Last two modes of operation enable us to test all control algorithms in any of standard high-level languages, like Visual Basic, C++, C#, Delphi etc. Usage of lower level increases burden on communication channel, as a result delays will increase too.

Conveyor control in on/off mode is controlled directly from Slave2 with available digital inputs and outputs. Advanced control can be done from Master, or by communicating between Slave1 and Slave3 (Variable Frequency Drive) through Master.

Mode that was used till now was on/off control from Slave1 with possibility of adjusting speed from Master.

Handshaking with packing machine is by two signals: with Product Ready combinatorial weigher informs packing machine that batch is ready and waiting, with Product Discharge packing machine requests batch and by reception of this signal DH is opened and product is released. Lower level functions are implemented and Master can control this process too.

3.3 Master Application Software

Software in Master side was developed in Visual Basic programming language with some native Win32 API calls. It contains two main software parts: HMI part that handles interaction with user and MODBUS Master protocol implementation. MODBUS protocol part was implemented as time triggered polled stateflow. Two freeware modules were used. Module crpTmr6 [19] was used for precise fast responsive time triggering. For serial port communication functions from CommIO [20] module were used.

When transmitting, RS485 driver must be enabled. For this function usually is used DTR line of serial port, but during tests it was found that delay in releasing driver was not tolerable and solution is found with auto switching RS485/RS232 interface. Clock for time triggering was 1000Hz (T=1ms). During one time tick the following sequence is executed:

1. Go to MODBUS Current State Entry.
2. If State Transition Flag is set execute Entry Action
3. Check, in a given order, if any of Transition Conditions is true
4. If true conditions is found then Exit Action of Current State is executed. Transition Action is executed. Current State changes to New State. State Transition Flag is set. Continue.
5. If no true condition is found then execute During Action. Continue.

If any of mentioned actions may or may not be present.

Only one state sequence per time tic is processed. HMI reaction to user inputs is processed in event generated by user action. It may result in local actions, or it may initiate sequence of local and network actions. Every request from slaves must be followed by response. Response itself may initiate another sequence of actions. Display content may change as result of actions.

For every HMI function group, screens were designed that give information in appropriate manner that is intuitive and informative [1] [2]. Also are given places for user interaction in form of buttons, sliders, check boxes, and numeric and text fields.

First screen of HMI after power-on, Fig. 4, gives us results software integrity test and results of hardware tests of integrated units. Fig.5 shows main operating screen. Only product related data are shown. Operating screen with detailed information of weights on each hopper, vibrating intensities for RF
and CF, and mean weight supplied by each head is presented in Fig. 6. Additional information is displayed for testing purpose.

### 3.4 Bus Transactions

Activity in bus has two modes of operation. First, default one, is organized in cyclic fashion. Time frame is divided by following sequence of transactions:

1. Slave1 Read/Write Multiple registers
2. Slave2 Read/Write Multiple registers
   - 3. Slave3 Read/Write Multiple registers
   - 4. Slave1 any request
   - 5. Slave2 any request

Time period of one frame is 0.1 s. Transactions 1 and 2, at beginning of frame, are transaction for real-time communication. They use function code 23 for writing and reading in same cycle. Transaction 3 to Slave3 (VFD) currently was used only for manual speed set. Transactions 4 and 5 are used for acyclic real-time communications or for non real-time communications, with priority to acyclic real-time.

Second mode of bus activity is based on shorter asynchronous transactions. Requests are divided into two classes of high and low priority level [21]. Master sends cyclically general request to slaves and slave sends data with highest priority.

Sequence of transactions for one pack cycle is:

1. Set Time/Intensity for Circular Feeder and start
2. Set Time/Intensity for Radial Feeder and start
3. Release product from HH to WH
4. Get weight data from Slave1 and display it
5. Inform Slave2 for weight ready on WH
6. Release product from WH to MH
7. Best combination is sent to Slave2
8. Release selected combination
9. Inform packing machine and wait for request
10. Request arrived and product is released

### 4. Initial Test Results

After all functions were implemented, system was tested in open-loop. Range of control values for CF and RF was found. Steady state characteristic for potato chips speed of travelling as function of RF intensity were constructed. And relation is in agreement with results in [22] for bulk transport on vibratory feeder. Offset type difference between RF was observed. Using intensity of RF as control variable was not effective. Time of operation was chosen as control variable for RF with intensity set to some optimal value. Number of MH for discharge was set to 3, as advised in [14]. With 75g and 175g packs of potato chips system worked almost without any adjustment. With smaller weights, 45g and 25g, stopped conditions increased and were quite frequent during 25g packing. It was observed that cycle time of RF and not working time of RF is important.
control parameter. Closed-loop algorithms are being tested. Initial test results for closed-loop control could reach continuous speed up to 57 packs/second for 25g packs. Same speed was reached when two opposite heads were switched-off. This suggests that speed limit is caused by volume of incoming product. Controlling variance and mean value for this multi-input multi-output system gave promising results.

5. Conclusion
Combinatorial weighers are in production lines that involve weighing with narrow tolerance is common and irreplaceable. Of importance is its operation without human intervention. Function of combinatorial weighing includes two processes: process of vibratory conveying and of combinatorial selection of best candidate from number of choices. To have the possibility of studying this process and test different control strategies experimental setup was created. For flexibility of implementation and for testing different control strategies, control over network communicating units is selected. Visual Basic with its open architecture and with built in functionality for user interface design is easily used for creating HMI applications. This was shown by designing this application down to the protocol implementation. Experimental setup has proved flexibility of MODBUS for implementing different modes of activity in bus and ability to handle real-time control for slow processes. Ability of personal computer (PC) of moderate performance to handle HMI, low level protocol functionality and real time control loops, including combinatorial calculation without any dedicated hardware, is proved. To enable flexibility of control, set of software functions were implemented on slave side. Open-loop tests for weighing product on real production line of potato chips were done. In closed-loop control continuous speed up to 57 packs/second for 25g packs could be reached. Variance and mean value control loops were used during closed-loop tests. Initial tests have shown flexibility of setup in implementing real-time closed-loop control over MODBUS communication channel for combinatorial weigher at potato chips production line.

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