Library for Systematic Search for Expressions

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Abstract: - In our previous work we showed that systematic search approach to inductive functional programming automation makes a remarkably efficient algorithm, but the applicability of the implemented program was limited by the very poor interpreter incorporated. In this paper, we present a library-based implementation of our systematic search algorithm which is supposed to be used with a well-known Haskell interpreter.

Key-Words: - inductive program synthesis, functional programming, library design, metaprogramming, Template Haskell, artificial intelligence

1 Introduction

MagicHaskeller is our inductive programming automation system by (exhaustively-and-efficiently-)generate-and-filter scheme for the lazy functional language Haskell. Until our previous work[1][2] it has been implemented as a stand alone program including a cheap interpreter for filtration of the generated programs. In this research we implemented its API library package which is supposed to be used with Glasgow Haskell Compiler interactive (GHCi) version 6.4 and higher.

The library implementation enjoys the following effects (some of which will be discussed further in Section 4):

• user-defined data types can be used;
• the full language features and libraries that GHC provides are supported for processing generated expressions;
• the usability when one wants to choose from different primitive sets has improved; for example, users can redefine reusable chunks of primitives such as list-related one, nat-related one, etc. and append them to obtain a suitable primitive set for the synthesis to be conducted.
• it has become easier to investigate the properties of the set of all the expressions that can be generated from the given type and the given primitive component set.

2 Background

2.1 MagicHaskeller

MagicHaskeller[1][2] is a project which emerged from our question: “Is genetic programming really more efficient than elaborated systematic exhaustive search for type-correct programs? If so, when and how is it? Has someone ever compared them?” In order to do the right things in the right order, in this project we first elaborate on optimized implementation of systematic exhaustive search for type correct programs, and then, if we become certain that we need heuristic approaches and that they really help improving the efficiency, we will build heuristic approaches on the top of our research on efficient systematic search algorithms under type constraints. Using techniques from software science such as monad for breadth-first search[3], generalized trie[4], monad transformer[5], and memoization using lazy data type[2], we have already shown that systematic exhaustive search for type-correct programs is a powerful alternative to genetic programming. The implemented program and library can be downloaded and tried from http://nautilus.cs.miyazaki-u.ac.jp/~skata/MagicHaskeller.html.

2.2 Template Haskell

Template Haskell (TH) [6] brings template metaprogramming to Haskell. In TH, expressions, sets of declarations, and types to be spliced have type \( \text{ExpQ} \), \( \text{Q [Dec]} \), and \( \text{TypeQ} \) respectively, where \( \text{ExpQ} = \text{Q Exp} \) and \( \text{TypeQ} = \text{Q Type} \), where monad Q is a wrapper to avoid name collision. A code block can be spliced by using \( \$\). An expression, a set of declarations, and a type can be quoted by using \( \ll| . | \rr\), \( \ll| . | \rr\), and \( \ll| . | \rr\) respectively.

For each of the primitive components we need the information of both of its semantics and its syntax, i.e., the value of the component which is used for interpretation and the name of the component which is used to print the generated expression. The most important role of TH in our library is to generate both of them without re-
3 Library specification

This section is devoted to the documentation of the implemented library. For each of the exported functions, its type and its short description are given.

3.1 Re-exported modules

This library implicitly re-exports the entities from `module Language.Haskell.TH as TH` and `module Data.Typeable from the Standard Hierarchical Library of Haskell`. Please refer to their documentations on types from them — in this documentation, types from TH are all qualified and the only type class used from module Typeable is Typeable.Typeable.

The following types are assigned to our internal data representations: `type Primitive` representing each primitive combinator and `type Memo` representing the memoization table.

3.2 Setting up the synthesis

Before synthesis, you have to define at least a memoization table (or you may define one once and reuse it for later syntheses). Other parameters are memoization depth and time out interval, which have default values. You may elect either to set those values to the ‘global variables’ using ‘set*’ functions, or hand them explicitly as parameters.

Note that the functions related to global variables are supposed to be used within the GHCi environment, and are not tested on other environments. They are implemented in the same way as GHCi implements global variables. For example, the global memoization table is implemented as:

```haskell
{-# NOINLINE refmemodeb #-
refmemodeb :: IORef Memo
refmemodeb = unsafePerformIO (newIORef defaultMD)
defaultMD = mkMemo []
```

Functions for creating memoization tables from the primitive components You can set your primitives like, e.g.,

```
s$ \langle + \rangle :: \text{Int} \to \text{Int} \to \text{Int}
\langle \text{a} \rangle, \lfloor \rfloor :: \lfloor \text{a} \rfloor
```

where the primitive set is consisted of \langle + \rangle specialized to type Int, \langle 0 \rangle specialized to type Int, \langle 'a' \rangle which has monomorphic type Char, and \lfloor \rfloor with polymorphic type \lfloor \text{a} \rfloor. As primitive components one can include any variables and constructors within the scope unless data types holding functions such as \lfloor \text{a} \to \text{b} \rfloor, \langle \text{Int} \to \text{Char}, \text{Bool} \rangle, etc., are involved. However, because currently ad hoc polymorphism is not supported by this library, you may not write

```
p :: TH.ExpQ \to TH.ExpQ
\text{type Primitive} = (\text{HValue}, TH.Exp, \text{String})
```

\(p\) is used to convert your primitive component set into its internal form. Usually its argument is the quasi-quote of a tuple of primitive components, but when it is not, singleton list will be generated.

The internal representation of each item of the primitive component set has type `Primitive = (HValue, TH.Exp, String)`, where `newtype HValue = HV (forall a.a)` is a universally quantified value representing the semantics of the primitive component, `TH.Exp` represents its syntactic aspect, and `String` is a string representation of its type. For example, \(p \langle((+)) :: Int \to Int, 0 :: Int, 'a' :: \lfloor a \rfloor\rangle\) reduces to

```
\langle HV (unsafeCoerce\langle ((+)) :: Int \to Int \rangle),
TH.VarE (TH.mkName "\(\langle +\rangle\)"),
"(((\to) Int) (((\to) Int) Int)))\rangle",
\langle HV (unsafeCoerce\langle (0 :: Int)\rangle),
TH.LitE (TH.IntegerL 0),
"Int")\rangle,
\langle HV (unsafeCoerce\langle 'a' \rangle),
TH.LitE (TH.CharL 'a'),
show (typeof 'a')\rangle,
\langle HV (unsafeCoerce\langle (\lfloor :: \forall a.a \to [a] \rfloor)\rangle),
TH.ConE (TH.mkName "[\rfloor"),
"([ [ [ 0 ]])\rangle
```

Use of `HValue` and `unsafeCoerce` for such purposes is taken from the source code of GHCi. The type consistency will not be assured statically by GHC but dynamically by our library. When the type is not given, it will be obtained via `typeof` method as is seen in the ‘\(\langle a\rangle\)’ case of the above example, although that will cause ambiguity error if the value is polymorphic.

```
setPrimitives :: [Primitive] \to \text{IO} ()
setPrimitives = setMemo.mkMemo
mkMemo :: [Primitive] \to \text{Memo}
setMemo :: Memo \to \text{IO} ()
```

Memoization depth

```
setDepth :: Int \to -- memoization depth \text{IO} ()
```

\(^1\)Here we stripped some of the qualifications by module names in order to avoid clutter.
Functions related to timeout Because the library generates all the expressions including those with non-linear recursions, you should note that there exist some expressions which take extraordinarily long time.[2] Imagine a function that takes an integer n and increments 0 for 2^n times. Another example that has smaller program size is blah in

<table>
<thead>
<tr>
<th>setDepth</th>
<th>can be used to set the memoization depth. Subexpressions within this size are memoized, while greater expressions will be recomputed (to save the heap space).</th>
</tr>
</thead>
<tbody>
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</tr>
</tbody>
</table>
| blih | Int → Int → Int  
| blih 0 i = i  
| blih k i = foo (blih (k - 1)) i  
| foo h 0 = 2  
| foo h s = h (f h (s - 1)) |

Giving small values in your examples is a good policy, but even then, time out is useful. For this reason, time out is by default taken after 1 second since each invocation of evaluation. This default behavior can be overridden by the following functions.

| setTimeout | :: Int → time in seconds IO ()  
| unsetTimeout | :: IO () |

Defining generator functions automatically In this case “automatically” does not mean “inductively” but “deductively using TH”.

| define | :: String → Integer → TH.ExpQ → TH.Q [TH.Dec]  
| define | eases use of this library by automating some function definitions. For example, $defn "Foo" 15 (p [[1 :: Int, (+) :: Int → Int → Int]])$ is equivalent to

| memoFoo | = mkMemo (p [[1 :: Int, (+) :: Int → Int → Int]])  
| everyFoo | = everything 15 memoFoo  
| filterFoo | :: Filter  
| filterFoo pred | = filterThen pred everyFoo |

where

| type | Everything = Typeable a ⇒ [[(TH.Exp,a)]]  
| type | Filter = 

| Typeable a ⇒ (a → Bool) → IO [[(TH.Exp,a)]] |

3.3 The synthesizers

Now is the time for defining the synthesizer functions. Most of the functions here just filter everything with the predicates you provide.

| quick starters |  
| findOne | :: Typeable a ⇒ (a → Bool) → TH.Exp  
| printOne | :: Typeable a ⇒ (a → Bool) → IO ()  
| printAny | :: Typeable a ⇒ (a → Bool) → IO ()  

| findOne pred | finds an expression e that satisfies pred e ≡ True, and returns it in TH.Exp. printOne prints the expression found first. printAny prints all the expressions satisfying the given predicate.  
| filterFirst | :: Typeable a ⇒ (a → Bool) → [[(TH.Exp,a)]]  
| filterThen | :: Typeable a ⇒ (a → Bool) → [[(TH.Exp,a)]] → IO [[(TH.Exp,a)]] |

filterFirst is like printAny, but by itself it does not print anything. Instead, it creates a stream of lists of expressions represented in tuples of TH.Exp and the expressions themselves. The stream contains expressions consisted of n primitive components at the nth element (n = 1, 2,...), and thus can be viewed as Spivey[3]’s Matrix data type.

| filterThen | may be used to further filter the results.  
| expression generators | These functions generate all the expressions that have the type you provide.  
| everything | :: Typeable a ⇒ 

| Int | → memoization depth.  
| Memo | → memo table  
| getEverything | :: Typeable a ⇒ IO [[(TH.Exp,a)]] |

everything generates all the expressions that fit the inferred type, and their representations in the TH.Exp form.

getEverything is like everything, but uses the global values set with set* functions.

| unifyable, matching | :: Int → Memo → TH.Type → [[TH.Exp]] |

These two functions are like everything, but take TH.Type as an argument, which may be polymorphic. For example, printQ (forall a.a → a → a) >>= return.unifiable 10 memo) will print all the expressions using memo whose types unify with forall a.a → a → a.
Pretty printer  These are utility functions for pretty printing the results:

\[
\begin{align*}
\text{pprs} & : \left(\left[\left(\text{TH.}\text{Exp}, a\right)\right]\right) \rightarrow \text{IO} () \\
\text{printQ} & : \text{Ppr} a \Rightarrow \text{Q} a \rightarrow \text{IO} ()
\end{align*}
\]

4 Use cases

This section presents many interesting results our library implementation has brought around. Unless otherwise specified, we assume that the Haskell source file in Table 1 is preloaded into GHCi.

User defined data types  From its birth our algorithm is designed to be able to deal with user-defined data types, but such ability had been suppressed by our poor interpreter in our previous work. Now that we can use the full power of GHCi, this potential ability is uncovered.

Suppose you have a Haskell source file shown in Table 2. You may synthesize programs by using it. Table 3 shows a sample interaction.

Actually the first two of the generated expressions implement different functions. You may want to use QuickCheck\[7\] for random testing in order to confirm this fact. You define \textbf{instance} \texttt{Arbitrary a ⇒ Arbitrary (Tree a)} and reload the source. Then, you filter further, as shown in Table 4. Of course, you may use QuickCheck to automatically filter the results further, though in this case you have to write a driver.

Using predicates in general other than examples  Unlike programming by example approaches, our library accepts any predicate as well as I/O example pairs. Table 5 shows an example of finding recursive form definitions of the Fibonacci function from its closed-form solution.

Managing your primitive set  Selecting your primitive set has also become easier. Access to your source file is required less frequently now, as shown in Table 6.

Browsing everything  Just browsing all the expressions having the given type is also interesting (Table 7).

It has become easier to tell some properties that the syntheses have. The example in Table 8 counts the number of expressions with each size. ²

²Note that these are not the ‘exact’ numbers of all the expressions, because our algorithm suppresses some of the duplicate expressions doing the same thing having different syntaxes, but still there remains such redundancy. Rather, you should interpret these numbers suggesting the ability of our algorithm.

5 Conclusions

We presented a library-based implementation of our systematic program generator algorithm which is supposed to be used with Glasgow Haskell Compiler interactive. By using several use cases we showed that this library expands the applicability of our algorithm.

Still, however, there are things to be done in order to enhance the usability of our algorithm:

- this algorithm can enhance Hoogle[8], that is a Haskell API search engine, by providing the ability to generate expressions;
- desirable set of primitive components (e.g. constructors and paramorphisms) for predefined types and their (better) pretty printers should be predefined, and those for user-defined types should be auto-generated;
- there could be a wrapper that automatically curry and uncurry tuple arguments

Also, support of type classes and further efficiency improvement could improve our algorithm.

References

Table 1: Default primitive component library

```haskell
{-# OPTIONS -fth #-}
module LibTH (module LibTH, module MagicHaskeller) where

import MagicHaskeller

initialize = do
    setPrimitives (list ++ nat ++ mb ++ bool ++ $(p \ (hd :: [a] → Maybe a, (+) :: Int → Int → Int)]))
    setDepth 15
    -- Specialized memoization tables. Choose one for quicker results.
    mlist :: Memo
    mlist = mkMemo (list ++ nat ++ mb ++ bool ++ $(p \ (hd :: [a] → Maybe a, (+) :: Int → Int → Int)]))
    mlistnat :: Memo
    mlistnat = mkMemo (list ++ nat ++ $(p \ ((+) :: Int → Int → Int)]))
    hd :: [a] → Maybe a
    hd [] = Nothing
    hd (x:_) = Just x
    mb = $(p \ (Nothing :: Maybe a, Just :: a → Maybe a, caseMaybe :: Maybe b → a → (b → a) → a)]
    caseMaybe :: Maybe b → a → (b → a) → a
    caseMaybe Nothing xf = x
    caseMaybe (Just y) xf = f y
    nat = $(p \ (0 :: Int, succ :: Int → Int, nat para :: Int → a → (Int → a → a) → a)]
    -- Nat paramorphism
    nat para :: Int → a → (Int → a → a) → a
    nat para 0 xf = x
    nat para (i + 1) xf = f (i + 1) xf
    list = $(p \ ([ :: [a], (::) :: a → [a] → [a], list para :: [b] → a → (b → [b] → a → a) → a)])
    -- List paramorphism
    list para :: [b] → a → (b → [b] → a → a) → a
    list para [] xf = x
    list para (ys:ys) xf = f (ys : list para ys xf)
    bool = $(p \ (True, False, iF :: Bool → a → a)]
    iF :: Bool → a → a
    iF True xf = t
    iF False xf = f
```

Table 2: Component library file with a `data` declaration

```haskell
import MagicHaskeller

data Tree a = Leaf a | Branch (Tree a) (Tree a) deriving (Eq, Typeable, Show)
tree para :: Tree a → (a → r) → (Tree a → tree para a → r → r → r) → r
tree para (Leaf x) f g = f x
tree para (Branch t u) f g = g t u (tree para t f g) (tree para u f g)
$\text{define "Tree" 10 (p \ [(Leaf :: a → Tree a, Branch :: Tree a → Tree a → Tree a,
  tree para :: Tree a → (a → r) → (Tree a → tree para a → r → r → r) → r)]})
Table 3: Sample interaction with a user-defined type

> ghci -v0 -fth -package MagicHaskeller TreeExample.hs

*Main> x <- filterTree (\f -> f (Branch (Branch (Leaf 1)) (Leaf 2)) (Leaf (3::Int))) == Branch
(Leaf 3) (Branch (Leaf 2)) (Leaf (1::Int)))

*Main> pprs x
\a -> tree_para a (\b -> Leaf b) (\b c d e -> Branch e d)
\a -> tree_para a (\b -> Leaf b) (\b c d e -> Branch c d)
\a -> tree_para a (\b c -> c) (\b c d e f -> Branch c (d (e b))) a
\a -> tree_para a (\b c -> c) (\b c d e f -> Branch c (d (d b))) a
\a -> tree_para a (\b c -> c) (\b c d e f -> Branch (e c)) (d b) a
\a -> tree_para a (\b c -> c) (\b c d e f -> Branch (e (e c))) (d b) a
Interrupted.

*Main>

Table 4: Filtering further based on the results from QuickCheck

*Main> :reload

*Main> quickCheck (\a -> tree_para a (\b -> Leaf (b::Int)) (\b c d e -> Branch e d) ==
tree_para a (\b -> Leaf (b::Int)) (\b c d e -> Branch c d))
Falsifiable, after 29 tests:
Branch (Leaf (-1)) (Leaf (2)) (Leaf (2)) \$ Branch (Leaf (-2)) (Leaf (2))
Branch (Leaf (-2)) (Leaf (2)) (Leaf (2)) \$ Branch (Leaf (-1)) (Leaf (2))
Branch (Leaf (2)) (Leaf (-2)) (Leaf (-1)) \$ Branch (Leaf (-1)) (Leaf (-1))
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Branch (Leaf (2)) (Leaf (2)) (Leaf (2)) \$ Branch (Leaf (-1)) (Leaf (-2)) (Leaf (2))
Interrupted.

*Main>

Table 5: Finding recursive definitions from closed-form solution

LibTH> initialize
LibTH> let phi = (1 + sqrt 5)/2
LibTH> printAny (\f -> all (\n -> (f :: Int->Int) n == round ((phi^n - (1-phi)^n) / sqrt 5)) [0..9])
\a -> nat_para a (\b c -> b) (\b c d e -> c e ((+) d e)) 0 (succ 0)
\a -> nat_para a (\b c -> b) (\b c d e -> c e ((+) e d)) 0 (succ 0)
\a -> nat_para a (\b c -> b) (\b c d e -> c ((+) d e) d) 0 (succ 0)
\a -> nat_para a (\b c -> b) (\b c d e -> c ((+) e d) d) 0 (succ 0)
\a -> nat_para a (\b c -> b) (\b c d e -> c (suc e) ((+) d e)) 0 0
\a -> nat_para a (\b c -> b) (\b c d e -> c (suc e) ((+) e d)) 0 0
Interrupted.
Table 6: Selecting your primitive set

(example cont’d from Table 5)

*LibTH> -- elect to use specialized primitive set to make your synthesis quicker
*LibTH> setPrimitives (nat ++ $(p [| (hd :: [a] -> Maybe a, (+) :: Int -> Int -> Int) |]))
*LibTH> printAny (\f -> all (\n -> (f :: Int->Int) n == round ((phi^n - (1-phi)^n) / sqrt 5) ) [0..9])
  \a -> nat_para a (\b c -> b) (\b c d e -> c e ((+) d e)) 0 (succ 0)
  \a -> nat_para a (\b c -> b) (\b c d e -> c e ((+) e d)) 0 (succ 0)
  \a -> nat_para a (\b c -> b) (\b c d e -> c ((+) e d) d) 0 (succ 0)
  \a -> nat_para a (\b c -> b) (\b c d e -> c (succ e) ((+) d e)) 0 0
  \a -> nat_para a (\b c -> b) (\b c d e -> c (succ e) ((+) e d)) 0 0
Interrupted.

Table 7: Browsing all the expressions

*LibTH> initialize
*LibTH> printQ (\[t| forall a. a->a->a |] >>= return . unifyable 10 mall)
  \a b -> b
  \a b -> a
  \a b -> 0
  \a b -> True
  \a b -> False
  \a b -> []
  \a b -> Nothing
  \a b -> succ 0
  \a b -> succ b
  \a b -> succ a
  \a b -> hd []
  \a b -> Just 0
  \a b -> Just True
  \a b -> Just False
  \a b -> Just []
  \a b -> Just Nothing
  \a b -> succ (succ 0)
  \a b -> succ (succ b)
  \a b -> succ (succ a)
Interrupted.

Table 8: Counting the number of expressions with each size

*LibTH> initialize
*LibTH> runQ \[t| forall a. a->a->a |] >>= print . map length . unifyable 10 mall
[7,9,34,76,223,651,2249,8053,31169,Interrupted.